

On Generalized Entropy Measures and Non-extensive Statistical Mechanics

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Abstract

Product probability property, known in the literature as statistical independence, is examined first. Then various generalized entropies are introduced, all of which give generalizations to Shannon entropy. It is shown that the nature of the recursivity postulate automatically determines the logarithmic functional form for Shannon entropy. Due to the logarithmic nature, Shannon entropy naturally gives rise to additivity, when applied to situations having product probability property. It is argued that the natural process is non-additivity leading to non-extensive statistical mechanics, even in product probability property situations and additivity can hold due to the involvement of a recursivity postulate leading to a logarithmic function. Various generalizations, including Mathai's generalized entropy are introduced and some of the properties are examined. Various situations are examined where Mathai's entropy leads to pathway models, Tsallis statistics, superstatistics, power law and related differential equations. Connection of Mathai's entropy to Kerridge's measure of "inaccuracy" is also explored.

Key words: Generalized entropies, Mathai's entropy, characterizations, additivity, non-extensivity, Tsallis statistics, superstatistics, power law.

AMS Subject Classifications:

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1. Introduction

Extensive or non-extensive statistical mechanics arise from the additivity or non-additivity of the corresponding entropy measures. Mathai and Rathie (1975,1976) consider various generalizations of Shannon entropy [Shannon(1948)], called entropies of order α , and give various properties, including additivity property, and characterization theorems. Tsallis (1988) considered a generalization of Shannon entropy, which is a variant of Havrda-Charvát (1967) entropy of order α with a different normalizing factor, and showed that his generalized entropy leads to power law, Tsallis statistics and to non-additivity thereby leading to non-extensive statistical mechanics. Non-extensive statistical mechanics is a hot topic now and a large number of papers have been published on the topic of Tsallis statistics and non-extensive statistical mechanics. In 2006 Mathai introduced a generalized entropy of order α , which is connected to a measure of uncertainty in a probability scheme, Keridge's (1961) concept of inaccuracy in a scheme, and Mathai's pathway model. Apart from giving a characterization theorem it is shown that Tsallis statistics [Tsallis (1988)] and superstatistics [Beck(2006), Beck and Cohen (2003)] are available as special cases from there [Mathai and Haubold (2006)].

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Mathai's entropy $M_{k,\alpha}(P)$ is a non-additive entropy and his measure $M_{k,\alpha}^*(P)$ is an additive entropy. It is also shown that maximization of the continuous analogue of $M_{k,\alpha}(P)$, denoted by $M\alpha(f)$, gives rise to various functional forms for f , depending upon the types of constraints on f . It is shown that Tsallis statistics, superstatistics, pathway model and a large number of other models are available from an optimization of $M\alpha(f)$ under various types of constraints.

Too much emphasis is placed on the fact that Shannon entropy satisfies the additivity property, leading to extensivity. It will be shown that when the product probability property (PPP) holds then a logarithmic function can give a sum and a logarithmic function enters into Shannon entropy due to the assumption introduced through a certain type of recursivity postulate. The concept of statistical independence will be examined in Section 1 to illustrate that simply because of PPP one need not expect additivity to hold or that one should not expect this PPP should lead to extensivity. The types of non-extensivity, associated with the various generalized entropies, are pointed out even when PPP holds. The nature of non-extensivity that can be expected from a multivariate distribution, when PPP holds or when there is statistical independence of the random variables, is illustrated by taking a trivariate case.

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Maximum entropy principle is examined in Section 2. It is shown that optimization of various measures of entropies, in the continuous populations, under various types of constraints, leads to various types of models. It is shown that Mathai's entropy is the most convenient one to obtain various probability models including Tsallis statistics, superstatistics, pathway and other models.

Section 3 examines the types of differential equations satisfied by the various special cases of the pathway model.

1.1. Product probability property (PPP) or statistical independence of events

Let $P(A)$ denote the probability of the event A . If the definition $P(A \cap B) = P(A)P(B)$ is taken as the definition of independence of the events A and B then any event $A \in S$, and S the sure event are independent. But A is contained in S and then the definition of independence becomes inconsistent with the common man's vision of independence. Even if the trivial cases of the sure event S and the impossible event ϕ are deleted, still this definition becomes a resultant of some properties of positive numbers. Consider a sample space of n distinct elementary events. If symmetry in the outcomes is assumed then we will assign equal probabilities $\frac{1}{n}$ each to the elementary events. Let $C = A \cap B$. If A and B are independent then $P(C) = P(A)P(B)$. Let

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$$P(A) = \frac{x}{n}, P(B) = \frac{y}{n}, P(C) = \frac{z}{n}.$$

Then

$$\binom{x}{n} \binom{y}{n} = \binom{z}{n} \Rightarrow nz = xy, x, y, z = 1, 2, \dots, n-1, z < x, y \quad (1)$$

deleting S and ϕ . There is no solution for x, y, z for a large number of n , for example, $n = 3, 5, 7$. This means that there are no independent events in such cases and it sounds strange from a common man's point of view.

The term “independence” of events is a misnomer. This property should have been called product probability property or PPP of events or something like that. There is no reason to expect the information or entropy in a joint distribution to be the sum of the information contents of the marginal distributions when the PPP holds for the distributions, that is when the joint density or probability function is a product of the marginal densities or probability functions. We may expect a term due to the product probability to enter into the expression for the entropy in the joint distribution in such cases. But if the information or entropy is defined in terms of a logarithm, then naturally, logarithm of a product being the sum of logarithms, we can expect a sum coming in such situations. This is not due to independence or due to the PPP of the densities but due to the

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fact that a functional involving logarithm is taken thereby a product has become a sum. Hence it is unwise to put too much importance whether or not the entropy on the joint distribution becomes sum of the entropies on marginal distributions or additivity property when PPP holds.

1.2. How is logarithm coming in Shannon's entropy?

Several characterization theorems for Shannon entropy and its various generalizations are given in Mathai and Rathie (1975, 1976). Modified and refined versions of Shannon's own postulates are given as postulates for the first theorem characterizing Shannon entropy in Mathai and Rathie (1975). Apart from continuity, symmetry, zero-indifference and normalization postulates the main postulate in the theorem is a recursivity postulate, which in essence says that when the PPP holds then the entropy will be a weighted sum of the entropies, thus in effect, assuming a logarithmic functional form. The crucial postulate is stated here. Consider a multinomial population $P = (p_1, \dots, p_m)$, $p_i > 0$, $i = 1, \dots, m$, $p_1 + \dots + p_m = 1$, that is, $p_i = P(A_i)$, $i = 1, \dots, m$, $A_1 \cup \dots \cup A_m = S$, $A_i \cap A_j = \phi$, $i \neq j$. [If any p_i can take a zero value also then zero-indifferent postulate, namely that the entropy remains the same when an impossible event is incorporated into the scheme, is to be added]. Let $H_n(p_1, \dots, p_n)$ denote the entropy to be defined. Then the crucial recursivity postulate says that

$$H_n(p_1, \dots, p_{m-1}, p_m q_1, \dots, p_m q_{n-m+1}) = H_m(p_1, \dots, p_m) \quad (2)$$

$$+ p_m H_{n-m+1}(q_1, \dots, q_{n-m+1})$$

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$\sum_{i=1}^m p_i = 1$, $\sum_{i=1}^{n-m+1} q_i = 1$. This says that if the m -th event A_m is partitioned into independent events $P(A_m \cap B_j) = P(A_m)P(B_j) = p_m q_j$, $j = 1, \dots, n - m + 1$ so that $p_m = p_m q_1 + \dots + p_m q_{n-m+1}$ then the entropy $H_n(\cdot)$ becomes a weighted sum. Naturally, the result will be a logarithmic function for the measure of entropy.

There are several modifications to this crucial recursivity postulate. One suggested by Tverberg is that $n - m + 1 = 2$ and $q_1 = q, q_2 = 1 - q$, $0 < q < 1$ and $H_2(q, 1 - q)$ is assumed to be Lebesgue integrable in $0 \leq q \leq 1$. Again a characterization of Shannon entropy is obtained. In all the characterization theorems for Shannon entropy this recursivity property enters in one form or the other as a postulate, which in effect implies a logarithmic form for the entropy measure. Shannon entropy S_k has the following form:

$$S_k = -A \sum_{i=1}^k p_i \ln p_i, \quad p_i > 0, \quad i = 1, \dots, k, \quad p_1 + \dots + p_k = 1, \quad (3)$$

where A is a constant. If any p_i is assumed to be zero then $0 \ln 0$ is to be interpreted as zero. Since the constant A is present, logarithm can be taken to any base. Usually the logarithm is taken to the base 2 for ready application to binary systems. We will take logarithm to the base e .

1.3. Generalization of Shannon entropy

Consider again a multinomial population $P = (p_1, \dots, p_k)$, $p_i > 0, i = 1, \dots, k, p_1 + \dots + p_k = 1$. The following are some of the generalizations of Shannon entropy S_k .

$$R_{k,\alpha}(P) = \frac{\ln(\sum_{i=1}^k p_i^\alpha)}{1 - \alpha}, \quad \alpha \neq 1, \alpha > 0, \quad (4)$$

(Rényi entropy of order α of 1961)

$$H_{k,\alpha}(P) = \frac{\sum_{i=1}^k p_i^\alpha - 1}{2^{1-\alpha} - 1}, \quad \alpha \neq 1, \alpha > 0 \quad (5)$$

(Havrda-Charvát entropy of order α of 1967)

$$T_{k,\alpha}(P) = \frac{\sum_{i=1}^k p_i^\alpha - 1}{1 - \alpha}, \quad \alpha \neq 1, \alpha > 0 \quad \text{(Tsallis entropy of 1988)} \quad (6)$$

$$M_{k,\alpha}(P) = \frac{\sum_{i=1}^k p_i^{2-\alpha} - 1}{\alpha - 1}, \quad \alpha \neq 1, -\infty < \alpha < 2 \quad (7)$$

(Mathai entropy of 2006)

$$M_{k,\alpha}^*(P) = \frac{\ln(\sum_{i=1}^k p_i^{2-\alpha})}{\alpha - 1}, \quad \alpha \neq 1, -\infty < \alpha < 2, \quad (8)$$

(Mathai's additive entropy of 2006).

When $\alpha \rightarrow 1$ all the entropies of order α described above in (4) to (7) go to Shannon entropy S_k .

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$$\lim_{\alpha \rightarrow 1} R_{k,\alpha}(P) = \lim_{\alpha \rightarrow 1} H_{k,\alpha}(P) = \lim_{\alpha \rightarrow 1} T_{k,\alpha}(P) = \lim_{\alpha \rightarrow 1} M_{k,\alpha}(P) = \lim_{\alpha \rightarrow 1} M_{k,\alpha}^*(P) = S_k. \quad (9)$$

Hence all the above measures are called generalized entropies of order α .

Let us examine to see what happens to the above entropies in the case of a joint distribution. Let $p_{ij} > 0$, $i = 1, \dots, m, j = 1, \dots, n$ such that $\sum_{i=1}^m \sum_{j=1}^n p_{ij} = 1$. This is a bivariate situation of a discrete distribution. Then the entropy in the joint distribution, for example,

$$M_{m,n,\alpha}(P, Q) = \frac{\sum_{i=1}^m \sum_{j=1}^n p_{ij}^{2-\alpha} - 1}{\alpha - 1}. \quad (10)$$

If the PPP holds and if $p_{ij} = p_i q_j$, $p_1 + \dots + p_m = 1$, $q_1 + \dots + q_n = 1$, $p_i > 0$, $i = 1, \dots, m$, $q_j > 0$, $j = 1, \dots, n$ and if $P = (p_1, \dots, p_m)$, $Q = (q_1, \dots, q_n)$ then

$$\begin{aligned} (\alpha - 1)M_{m,\alpha}(P)M_{n,\alpha}(Q) &= \frac{1}{\alpha - 1} \left(\sum_{i=1}^m p_i^{2-\alpha} - 1 \right) \left(\sum_{j=1}^n q_j^{2-\alpha} - 1 \right) \\ &= \frac{1}{\alpha - 1} \left[\sum_{i=1}^m \sum_{j=1}^n p_i^{2-\alpha} q_j^{2-\alpha} - \sum_{i=1}^m p_i^{2-\alpha} - \sum_{j=1}^n q_j^{2-\alpha} + 1 \right] \\ &= M_{m,n,\alpha}(P, Q) - M_{m,\alpha}(P) - M_{n,\alpha}(Q). \end{aligned}$$

Therefore

$$M_{m,n,\alpha}(P, Q) = M_{m,\alpha}(P) + M_{n,\alpha}(Q) + (\alpha - 1)M_{m,\alpha}(P)M_{n,\alpha}(Q). \quad (11)$$

If any one of the above mentioned generalized entropies in (4) to (8) is written as $F_{m,n,\alpha}(P, Q)$ then we have the relation

$$F_{m,n,\alpha}(P, Q) = F_{m,\alpha}(P) + F_{n,\alpha}(Q) + a(\alpha)F_{m,\alpha}(P)F_{n,\alpha}(Q). \quad (12)$$

where

$$\begin{aligned} a(\alpha) &= 0 \quad (\text{Rényi entropy } R_{k,\alpha}(P)) \\ &= 2^{1-\alpha} - 1 \quad (\text{Havrda-Charvát entropy } H_{k,\alpha}(P)) \\ &= 1 - \alpha \quad (\text{Tsallis entropy } T_{k,\alpha}(P)) \\ &= \alpha - 1 \quad (\text{Mathai entropy } M_{k,\alpha}(P)) \\ &= 0 \quad (\text{Mathai entropy } M_{k,\alpha}^*(P)). \end{aligned} \quad (13)$$

When $a(\alpha) = 0$ the entropy is called additive and when $a(\alpha) \neq 0$ the entropy is called non-additive. As can be expected, when a logarithmic function is involved, as in the cases of $S_k(P)$, $R_{k,\alpha}(P)$, $M_{k,\alpha}^*(P)$, the entropy is additive and $a(\alpha) = 0$.

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1.4. Extensions to higher dimensional joint distributions

Consider a trivariate population or a trivariate discrete distribution $p_{ijk} > 0, i = 1, \dots, m, j = 1, \dots, n, k = 1, \dots, r$ such that $\sum_{i=1}^m \sum_{j=1}^n \sum_{k=1}^r p_{ijk} = 1$. If the PPP holds mutually, that is, pair-wise as well as jointly, which then will imply that

$$p_{ijk} = p_i q_j s_k, \quad \sum_{i=1}^m p_i = 1, \quad \sum_{j=1}^n q_j = 1, \quad \sum_{k=1}^r s_k = 1,$$
$$P = (p_1, \dots, p_m), \quad Q = (q_1, \dots, q_n), \quad S = (s_1, \dots, s_r).$$

Then proceeding as before, we have for any of the measures described above in (4) to (8), calling it $F(\cdot)$,

$$F_{m,n,r,\alpha}(P, Q, S) = F_{m,\alpha}(P) + F_{n,\alpha}(Q) + F_{r,\alpha}(S) + a(\alpha)[F_{m,\alpha}(P)F_{n,\alpha}(Q) + F_{m,\alpha}(P)F_{r,\alpha}(S) + F_{n,\alpha}(Q)F_{r,\alpha}(S)] + [a(\alpha)]^2 F_{m,\alpha}(P)F_{n,\alpha}(Q)F_{r,\alpha}(S) \quad (14)$$

where $a(\alpha)$ is the same as in (13). The same procedure can be extended to any multivariable situation. If $a(\alpha) = 0$ we may call the entropy additive and if $a(\alpha) \neq 0$ then the entropy is non-additive or we call the corresponding statistical mechanics extensive and non-extensive respectively.

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1.5. Crucial recursivity postulate

Consider the multinomial population $P = (p_1, \dots, p_k)$, $p_i > 0$, $i = 1, \dots, k$, $p_1 + \dots + p_k = 1$. Let the entropy measure to be determined through appropriate postulates be denoted by $H_k(P) = H_k(p_1, \dots, p_k)$. For $k = 2$ let

$$f(x) = H_2(x, 1 - x), \quad 0 \leq x \leq 1 \text{ or } x \in [0, 1]. \quad (15)$$

If another parameter α is to be involved in $H_2(x, 1 - x)$ then we will denote $f(x)$ by $f_\alpha(x)$. From (5) to (7) it can be seen that the generalized entropies of order α of Havrda-Charvát, Tsallis and Mathai, and Shannon entropy satisfy the functional equation

$$f_\alpha(x) + b_\alpha(x) f_\alpha\left(\frac{y}{1-x}\right) = f_\alpha(y) + b_\alpha(y) f_\alpha\left(\frac{x}{1-y}\right) \quad (16)$$

for $x, y \in [0, 1)$ with $x + y \in [0, 1]$, with the boundary condition

$$f_\alpha(0) = f_\alpha(1) \quad (17)$$

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$$\begin{aligned}
b_\alpha(x) &= 1 - x \text{ (Shannon entropy } S_k(P)) \\
&= (1 - x)^\alpha \text{ (Harvda-Charvát entropy } H_{k,\alpha}(P)) \\
&= (1 - x)^\alpha \text{ (Tsallis entropy } T_{k,\alpha}(P)) \\
&= (1 - x)^{2-\alpha} \text{ (Mathai entropy } M_{k,\alpha}(P)). \tag{18}
\end{aligned}$$

Observe that the normalizing constant at $x = \frac{1}{2}$ is equal to 1 for $H_{k,\alpha}(P)$ and it is different for other entropies. Thus equations (6),(7),(8), with the appropriate normalizing constants $f_\alpha(\frac{1}{2})$, can give characterization theorems for the various entropy measures. The form of $b_\alpha(x)$ is coming from the crucial recursivity postulate, assumed as a desirable property for the measures.

1.6. Continuous analogues

In the continuous case let $f(x)$ be the density function of a real random variable x . Then the various entropy measures, corresponding to the ones in (4) to (8) are the following:

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$$R_\alpha(f) = \frac{1}{1-\alpha} \ln \left[\int_{-\infty}^{\infty} [f(x)]^\alpha dx \right], \alpha \neq 1, \alpha > 0 \quad (19)$$

(Rényi entropy of order α)

$$H_\alpha(f) = \frac{1}{2^{1-\alpha} - 1} \left[\int_{-\infty}^{\infty} [f(x)]^\alpha dx - 1 \right], \alpha \neq 1, \alpha > 0 \quad (20)$$

(Havrda-Charvát entropy of order α)

$$T_\alpha(f) = \frac{1}{1-\alpha} \left[\int_{-\infty}^{\infty} [f(x)]^\alpha dx - 1 \right], \alpha \neq 1, \alpha > 0, \quad (21)$$

(Tsallis entropy of order α)

$$M_\alpha(f) = \frac{1}{\alpha-1} \left[\int_{-\infty}^{\infty} [f(x)]^{2-\alpha} dx - 1 \right], \alpha \neq 1, \alpha < 2 \quad (22)$$

(Mathai entropy of order α)

$$M_\alpha^*(f) = \frac{1}{\alpha-1} \ln \left[\int_{-\infty}^{\infty} [f(x)]^{2-\alpha} dx \right], \alpha \neq 1, \alpha < 2 \quad (23)$$

(Mathai's additive entropy of order α).

As expected, Shannon entropy in this case is given by

$$S(f) = -A \int_{-\infty}^{\infty} f(x) \ln f(x) dx \quad (24)$$

where A is a constant.

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Note that when PPP (product probability property) or statistical independence holds then in the continuous case also we have the property in (12) and (14) and then non-additivity holds for the measures analogous to the ones in (3),(5),(6),(7) with $a(\alpha)$ remaining the same. Since the steps are parallel a separate derivation is not given here.

2. Maximum Entropy Principle

If we have a multinomial population $P = (p_1, \dots, p_k)$, $p_i > 0, i = 1, \dots, k, p_1 + \dots + p_k = 1$ or the scheme $P(A_i) = p_i, A_1 \cup \dots \cup A_k = S, P(S) = 1, A_i \cap A_j = \phi, i \neq j$ then we know that the maximum uncertainty in the scheme or the minimum information from the scheme is obtained when we cannot give any preference to the occurrence of any particular event or when the events are equally likely or when $p_1 = p_2 = \dots = p_k = \frac{1}{k}$. In this case, Shannon entropy becomes,

$$S_k(P) = S_k\left(\frac{1}{k}, \dots, \frac{1}{k}\right) = -A \sum_{i=1}^k \frac{1}{k} \ln \frac{1}{k} = A \ln k \quad (25)$$

and this is the maximum uncertainty or maximum Shannon entropy in this scheme. If the arbitrary functional f is to be fixed by maximizing the entropy then in (19) to (21) we have to optimize $\int_{-\infty}^{\infty} [f(x)]^\alpha dx$ for fixed α , over all functional f , subject to the condition $\int_{-\infty}^{\infty} f(x) dx = 1$

and $f(x) \geq 0$ for all x . For applying calculus of variation procedure we consider the functional

$$U = [f(x)]^\alpha - \lambda[f(x)]$$

where λ is a Lagrangian multiplier. Then the Euler equation is the following:

$$\frac{\partial U}{\partial f} = 0 \Rightarrow \alpha f^{\alpha-1} - \lambda = 0 \Rightarrow f = \left(\frac{\lambda}{\alpha}\right)^{\frac{1}{\alpha-1}} = \text{constant}. \quad (26)$$

Hence f is the uniform density in this case, analogous to the equally likely situation in the multinomial case. If the first moment $E(x) = \int_{-\infty}^{\infty} xf(x)dx$ is assumed to be a given quantity for all functional f then U will become the following for (19) to (21).

$$U = [f(x)]^\alpha - \lambda_1[f(x)] - \lambda_2xf(x)$$

and the Euler equation leads to the power law. That is,

$$\frac{\partial U}{\partial f} = 0 \Rightarrow \alpha f^{\alpha-1} - \lambda_1 - \lambda_2x = 0 \Rightarrow f = c_1 \left[1 + \frac{\lambda_2}{\lambda_1}x\right]^{\frac{1}{\alpha-1}}. \quad (27)$$

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By selecting $c_1, \lambda_1, \lambda_2$ appropriately we can create a density out of (27). For $\alpha > 1$ and $\frac{\lambda_2}{\lambda_1} > 0$ the right side in (27) increases exponentially. If $\alpha = q > 1$ and $\frac{\lambda_2}{\lambda_1} = q - 1$ then we have Tsallis' q -exponential function from the right side of (27). If $\alpha > 1$ and $\frac{\lambda_2}{\lambda_1} = -(\alpha - 1)$ then (27) can produce a density in the category of a type-1 beta. From (27) it is seen that the form of the entropies of Havrda-Charvát $H_{k,\alpha}(P)$ and Tsallis $T_{k,\alpha}(P)$ are not convenient to produce densities. Hence Tsallis and others have considered a different constraint on $E(x)$. If the density $f(x)$ is replaced by its escort density, namely, $\mu[f(x)]^\alpha$ where $\mu^{-1} = \int_x [f(x)]^\alpha dx$ and if the expected value of x in this escort density is assumed to be fixed for all functional f then the U of (26) becomes

$$\begin{aligned}
 U &= f^\alpha - \lambda_1 f + \mu \lambda_2 x f^\alpha \\
 \text{and} \\
 \frac{\partial U}{\partial f} &= 0 \Rightarrow \alpha f^{\alpha-1} [1 + \mu \lambda_2 x] = \lambda_1 \Rightarrow f = \frac{\lambda^*}{(1 + \lambda_3 x)^{\frac{1}{\alpha-1}}} \Rightarrow \\
 f &= \lambda_1^* [1 + \lambda_3 x]^{-\frac{1}{\alpha-1}}
 \end{aligned}$$

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where λ_3 is a constant and λ_1^* is the normalizing constant. If λ_3 is take as $\lambda_3 = \alpha - 1$ then

$$f = \lambda_1^*[1 + (\alpha - 1)x]^{-\frac{1}{\alpha-1}}. \quad (28)$$

Then (28) for $\alpha > 1$ is Tsallis statistics for $\alpha > 1$. Then for $\alpha < 1$ also by writing $\alpha - 1 = -(1 - \alpha)$ one gets the case of Tsallis statistics for $\alpha < 1$. These modifications and the consideration of escort distribution are not necessary if we take Mathai's entropy. If we taken Mathai entropy $M_\alpha(f)$ and if we assume that the first moment in $f(x)$ itself is fixed for all functional f then the Euler equation gives

$$(2 - \alpha)f^{1-\alpha} - \lambda_1 + \lambda_2x = 0 \Rightarrow f = \bar{\lambda} \left[1 - \frac{\lambda_2}{\lambda_1}x \right]^{\frac{1}{1-\alpha}}$$

and for $\frac{\lambda_2}{\lambda_1} = 1 - \alpha$ we have Tsallis statistics

$$f = \bar{\lambda}[1 - (1 - \alpha)x]^{\frac{1}{1-\alpha}} \quad (29)$$

coming directly, where $\bar{\lambda}$ is the normalizing constant.

Let us start with Mathai entropy $M_\alpha(f)$ of (20) under the assumptions that $f(x) \geq 0$ for all x , $\int_a^b f(x)dx = 1$, $\int_a^b x^\delta f(x)dx$ is fixed for all functional f and for a specified $\delta > 0$, $f(a)$ is the same for all functional f , $f(b)$ is the same for all functional f , for some limits a and b , then the Euler equation becomes

$$(2 - \alpha)f^{1-\alpha} - \lambda_1 - \lambda_2x^\delta = 0 \Rightarrow f = c_1[1 + c_1^*x^\delta]^{\frac{1}{1-\alpha}}. \quad (30)$$

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If c_1^* is written as $-s(1 - \alpha)$, $s > 0$ then we have, writing f_1 for f ,

$$f_1 = c_1[1 - s(1 - \alpha)x^\delta]^{-\frac{1}{1-\alpha}}, \quad \delta > 0, \alpha < 1, \quad 0 \leq x \leq \frac{1}{[s(1 - \alpha)]^{\frac{1}{\delta}}} \quad (31)$$

where $1 - s(1 - \alpha)x^\delta > 0$. For $\alpha < 1$ or $-\infty < \alpha < 1$ the right side of (31) remains as a generalized type-1 beta model with the corresponding normalizing constant c_1 . For $\alpha > 1$, writing $1 - \alpha = -(\alpha - 1)$ the model in (31) goes to a generalized type-2 beta form, namely,

$$f_2 = c_2[1 + s(\alpha - 1)x^\delta]^{-\frac{1}{\alpha-1}}. \quad (32)$$

When $\alpha \rightarrow 1$ in (31) or in (32) we have an extended or stretched exponential form,

$$f_3 = c_3e^{-s x^\delta}. \quad (33)$$

If c_1^* in (30) is taken as positive then (30) for $\alpha < 1$, $\alpha > 1$, $\alpha \rightarrow 1$ will be increasing exponentially. Hence all possible forms are available from (30). The model in (31) is a special case of Mathai's pathway model and for a discussion of the matrix-variate pathway model see Mathai (2005). Special cases of (31) and (32) for $\delta = 1$ are Tsallis statistics.

Instead of optimizing $M_\alpha(f)$ of (22) under the conditions that $f(x) \geq 0$ for all x , $\int_a^b f(x)dx = 1$ and $\int_a^b x^\delta f(x)dx$ is fixed, let us optimize under the

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following conditions: $f(x) \geq 0$ for all x , $\int_a^b f(x)dx < \infty$ and the following two moment-like expressions are fixed quantities for all functional f ,

$$\int_a^b x^{(\gamma-1)(1-\alpha)} f(x)dx = \text{fixed}, \int_a^b x^{(\gamma-1)(1-\alpha)+\delta} f(x)dx = \text{fixed}.$$

Then the Euler equation becomes

$$(2 - \alpha)f^{1-\alpha} - \lambda_1 x^{(\gamma-1)(1-\alpha)} - \lambda_2 x^{(\gamma-1)(1-\alpha)+\delta} = 0 \Rightarrow$$

$$f = c x^{\gamma-1} [1 + c^* x^\delta]^{\frac{1}{1-\alpha}}$$

and for $c^* = -s(1 - \alpha)$, $s > 0$, we have Mathai's pathway model for the real scalar case, namely

$$f(x) = c x^{\gamma-1} [1 - s(1 - \alpha)x^\delta]^{\frac{1}{1-\alpha}}, \delta > 0, s > 0 \quad (34)$$

where c is the normalizing constant. For $\alpha < 1$, (34) gives a generalized type-1 beta form, for $\alpha > 1$ it gives a generalized type-2 beta form and for $\alpha \rightarrow 1$ we have a generalized gamma form. For $\alpha > 1$, (34) gives the superstatistics of Beck (2006) and Beck and Cohen (2003). For $\gamma = 1, \delta = 1$, (34) gives Tsallis statistics. All types of densities appearing in astrophysics problems and other areas are seen to be special cases of (34), a discussion of which may be seen from Mathai and Haubold (2006a). For example, (34) for $\delta = 2, \gamma = 3, \alpha \rightarrow 1, x > 0$ is the Maxwell-Boltzmann

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density; for $\delta = 2, \gamma = 1, \alpha \rightarrow 1, -\infty < x < \infty$ is the Gaussian density; for $\gamma = \delta, \alpha \rightarrow 1$ is the Weibull density. For $\gamma = 1, \delta = 2, 1 < q < 3$ we have the Wigner function $W(p)$ giving the atomic moment distribution in the framework of Fokker-Planck equation, see Douglas, et al (2006) where

$$W(p) = z_q^{-1} [1 - \beta(1 - q)p^2]^{\frac{1}{1-q}}, \quad 1 < q < 3. \quad (35)$$

Before closing this section we may observe one more property for $M_\alpha(f)$. As an expected value

$$M_\alpha(f) = \frac{1}{\alpha - 1} [E[f(x)]^{1-\alpha} - 1]. \quad (36)$$

But Kerridge's (Kerridge, 1961) measure of "inaccuracy" in assigning $q(x)$ for the true density $f(x)$, in the generalized form is

$$H_\alpha(f : q) = \frac{1}{(2^{1-\alpha} - 1)} [E[q(x)]^{\alpha-1} - 1], \quad (37)$$

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which is also connected to the measure of directed divergence between $q(x)$ and $f(x)$. In (37) the normalizing constant is $2^{1-\alpha} - 1$, the same thing appearing in Havrda-Charvát entropy. With different normalizing constants, as seen before, (36) and (37) have the same forms as an expected value with $q(x)$ replaced by $f(x)$ in (36). Hence $M_\alpha(f)$ can also be looked upon as a type of directed divergence or “inaccuracy” measure.

3. Differential Equations

The functional part in (34), for a more general exponent, namely

$$g(x) = \frac{f(x)}{c} = x^{\gamma-1} [1 - s(1 - \alpha)x^\delta]^{\frac{\beta}{1-\alpha}}, \quad \alpha \neq 1, \delta > 0, \beta > 0, s > 0 \quad (38)$$

is seen to satisfy the following differential equation for $\gamma \neq 1$.

$$\begin{aligned} x \frac{d}{dx} g(x) &= (\gamma - 1)x^{\gamma-1} [1 - s(1 - \alpha)x^\delta]^{\frac{\beta}{1-\alpha}} \\ &\quad - s\beta\delta x^{\delta+\gamma-1} [1 - s(1 - \alpha)x^\delta]^{\frac{\beta}{1-\alpha} [1 - \frac{(1-\alpha)}{\beta}]}. \end{aligned} \quad (39)$$

Then for $\delta = \frac{(\gamma-1)(\alpha-1)}{\beta}$, $\gamma \neq 1$, $\alpha > 1$ we have

$$x \frac{d}{dx} g(x) = (\gamma - 1)g(x) - s\beta\delta [g(x)]^{1 - \frac{(1-\alpha)}{\beta}} \quad (40)$$

$$\begin{aligned} &= (\gamma - 1)g(x) - s\delta [g(x)]^\alpha \text{ for } \beta = 1, \quad (41) \\ &\quad \gamma \neq 1, \delta = (\gamma - 1)(\alpha - 1), \alpha > 1. \end{aligned}$$

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For $\gamma = 1, \delta = 1$ in (38) we have

$$\frac{d}{dx}g(x) = -s[g(x)]^\eta, \quad \eta = 1 - \frac{(1 - \alpha)}{\beta} \quad (42)$$

$$= -s[g(x)]^\alpha \text{ for } \beta = 1. \quad (43)$$

Here (43) is the power law coming from Tsallis statistics.

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